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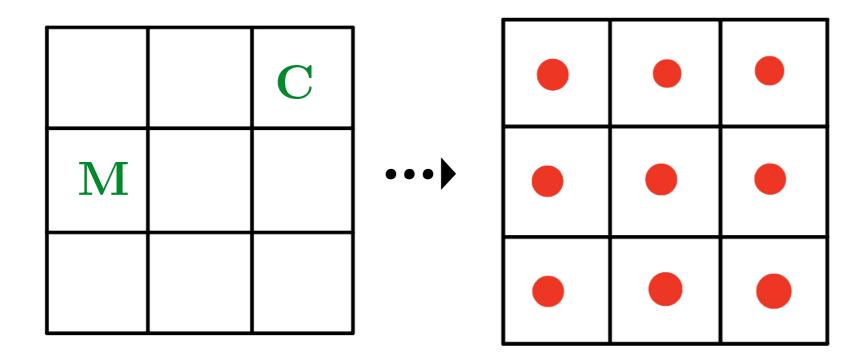
Main Question. Who has a winning strategy?

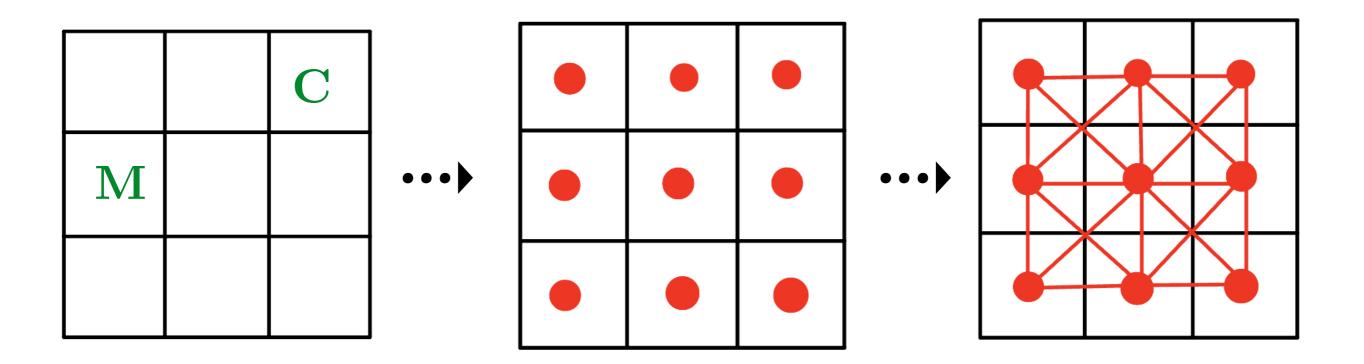
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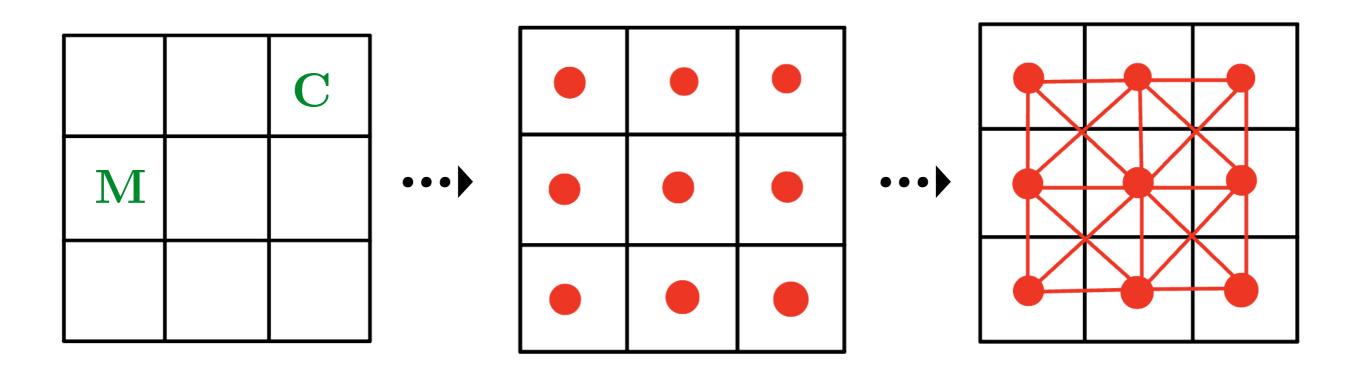
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Variation. Game begins by cat and mouse each choosing their starting position.

	\mathbf{C}
\mathbf{M}	

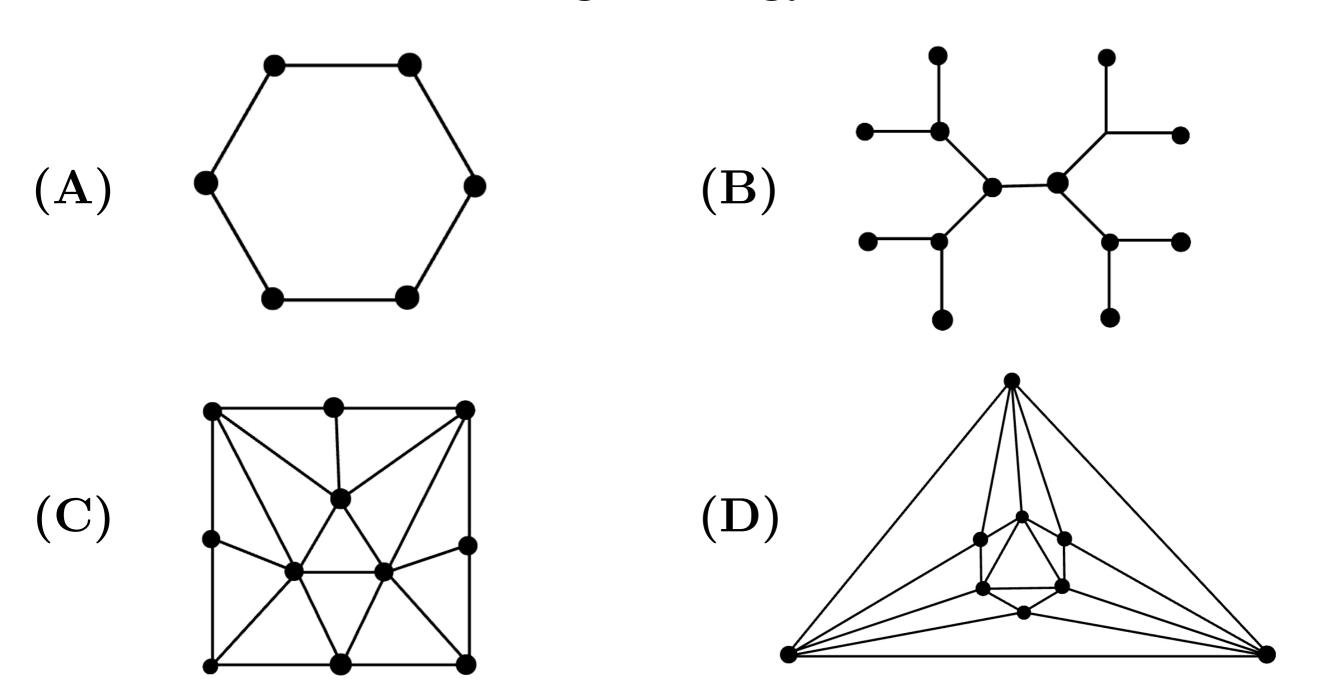




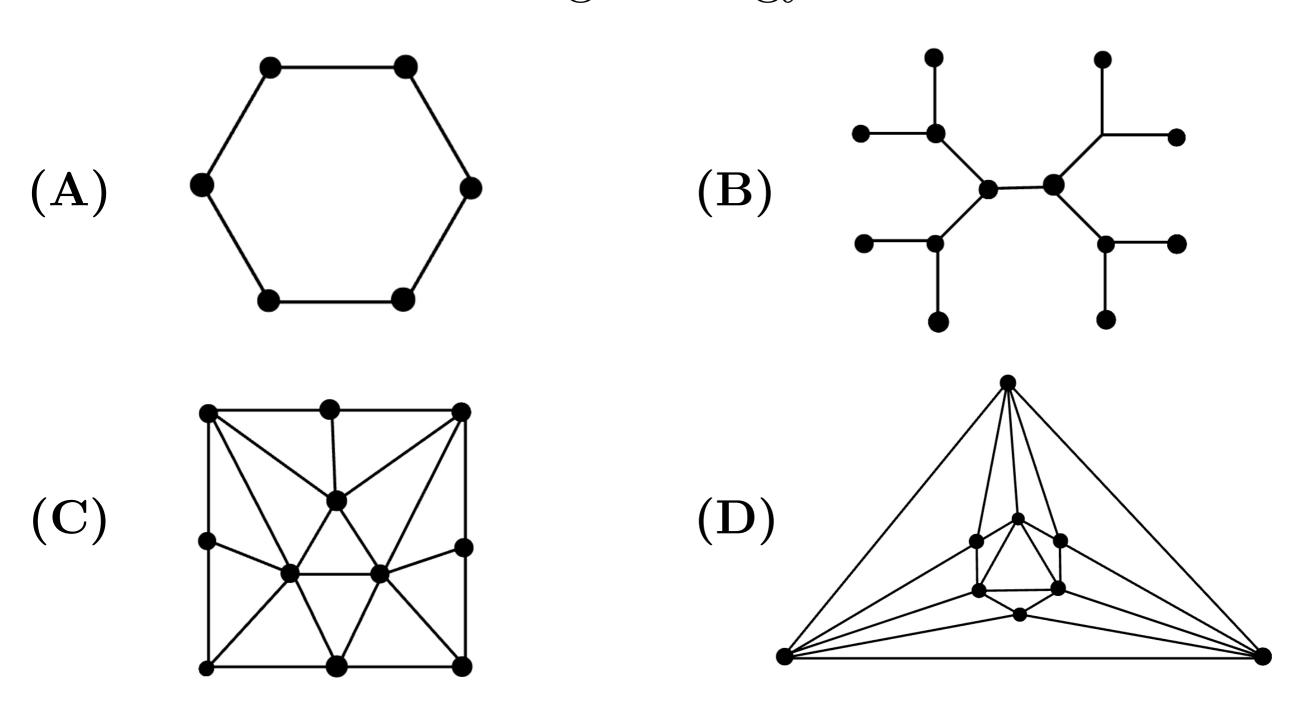


Cat and mouse play on vertices of the graph, moving to adjacent vertices.

Question. For these graphs, does either the cat or mouse have a winning strategy?



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How generally to tell if a graph is cat win?

Game theory: mathematical theory that analyzes strategy and decision making.

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Zermelo's theorem. In any* finite 2-player game without chance (e.g. chess, nim, cat-mouse) one player has a winning strategy.

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Zermelo's theorem. In any* finite 2-player game without chance (e.g. chess, nim, cat-mouse) one player has a winning strategy.

Proof is non-constructive!

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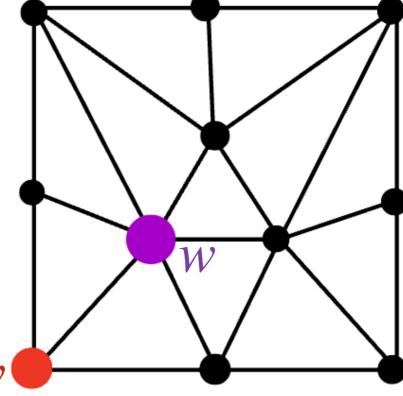
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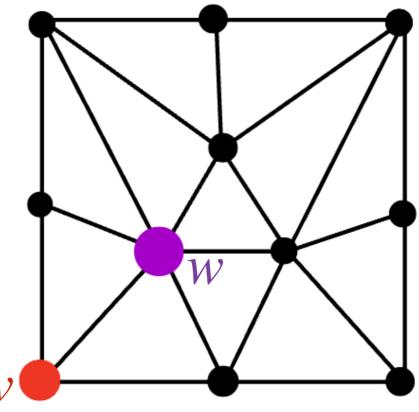
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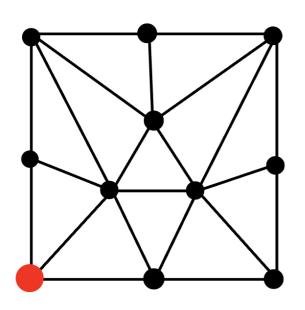
 $\underline{\text{dominated}}$ by w.

Remark. Having a dominated vertex is a local property of a graph.

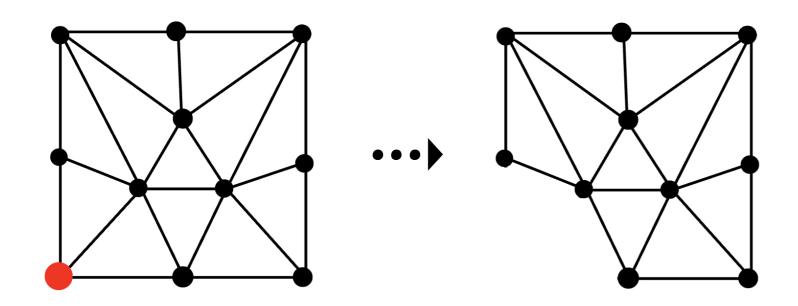


Fact. G graph. Assume v is dominated. If cat wins on $G \setminus v$, then cat wins on G.

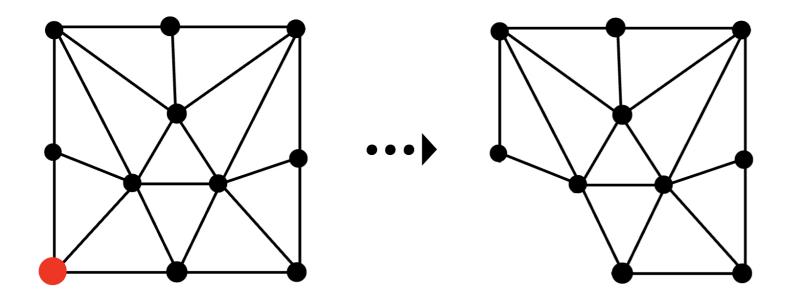
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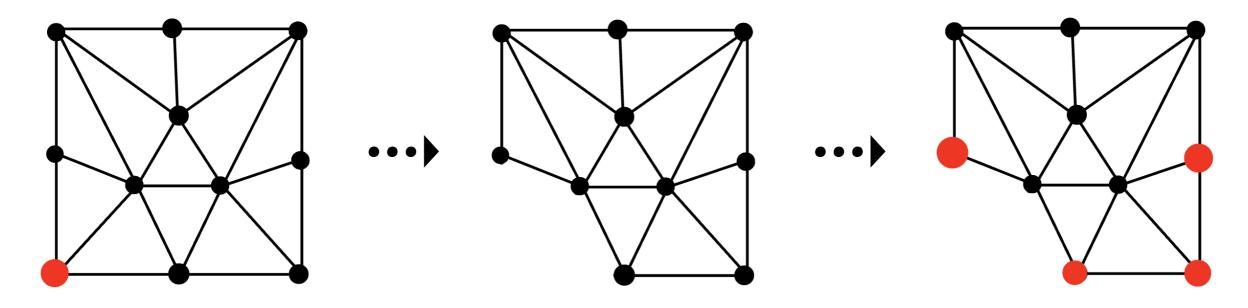
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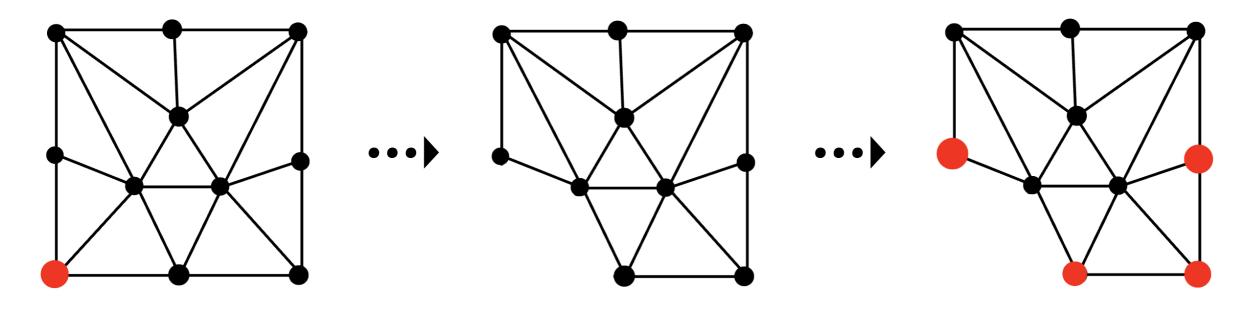
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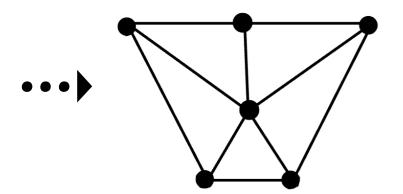


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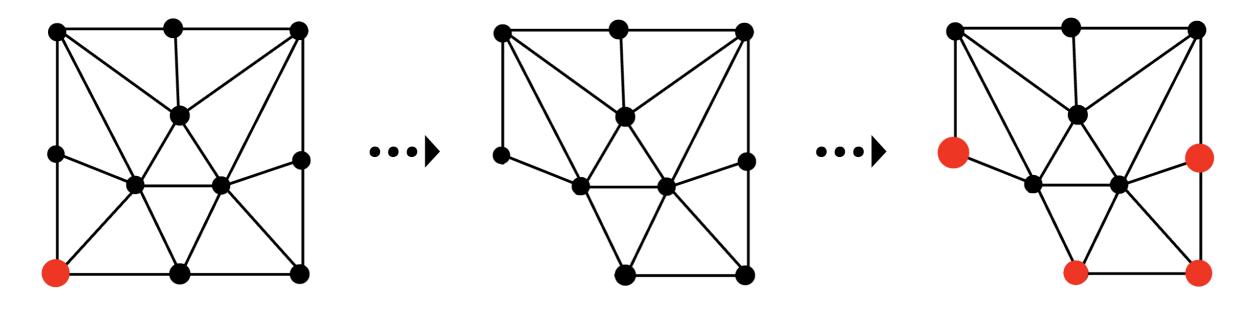


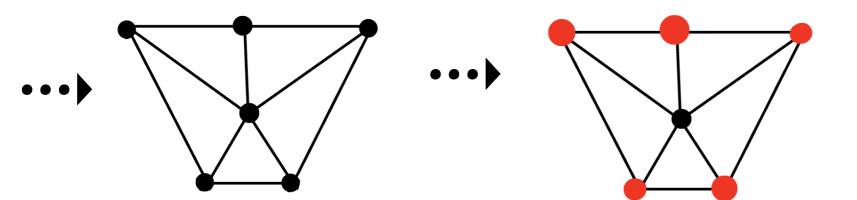
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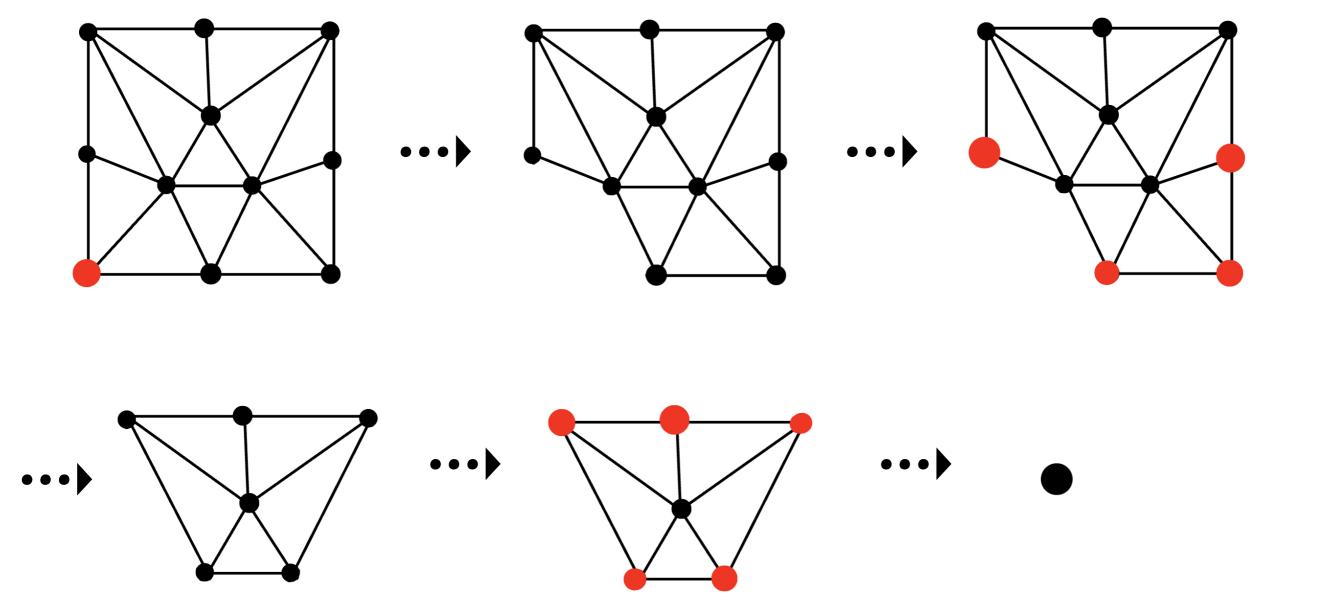


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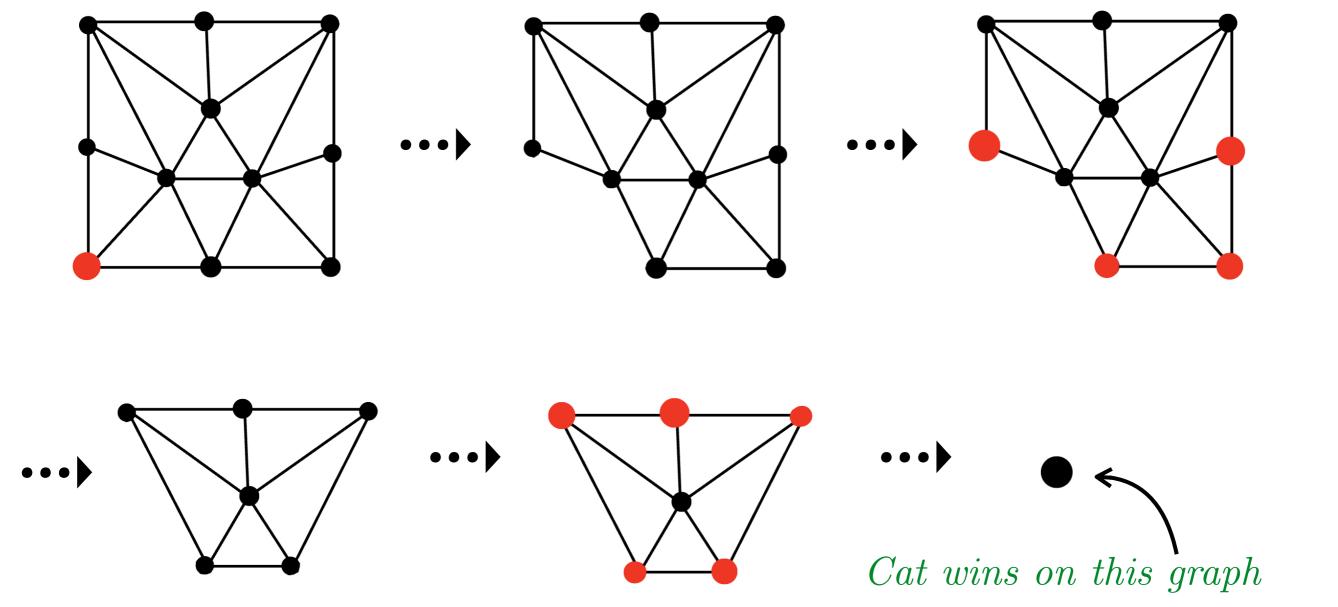
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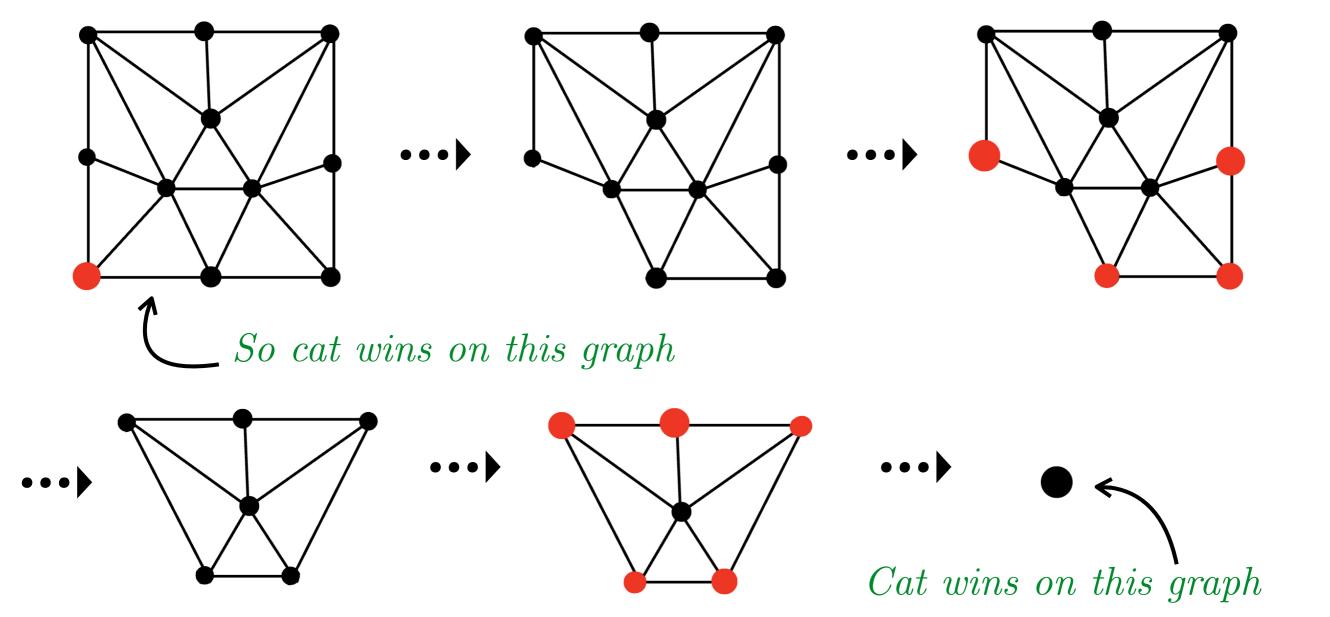
Using this, we have an inductive procedure to show cat has winning strategy.



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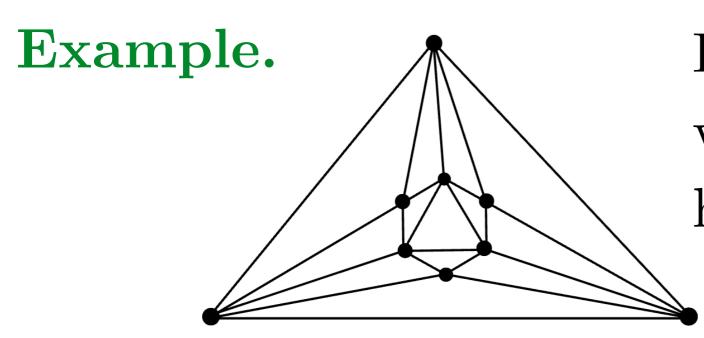
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Cat has winning strategy on G if and only if can order vertices $v_1, ..., v_n$ so that v_i is dominated by one of $v_{i+1}, ..., v_n$ for each i = 1, ..., n-1.

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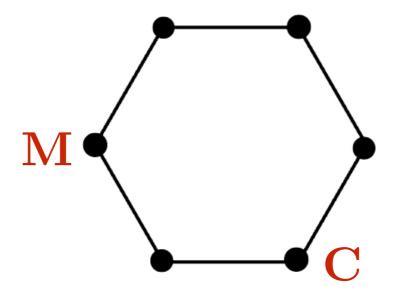
If G has no dominated vertex, then the mouse has winning strategy.

Variation: More cats!

How many cats are needed to ensure the cats win?

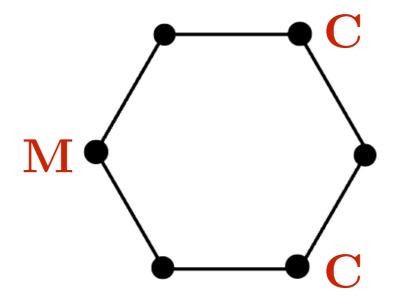
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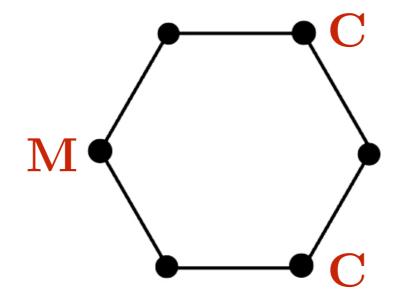
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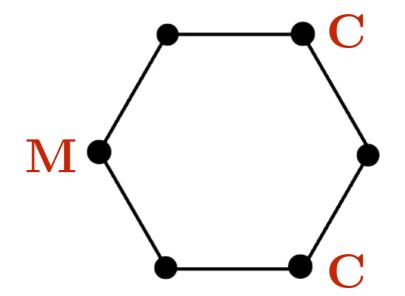
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Conjecture (Meyniel, 1985). In a graph with n vertices, don't need more than \sqrt{n} cats to catch a mouse. More precisely the maximum cat number among n-vertex graphs is $O(\sqrt{n})$.

